

Dynamic Programming (Contd).

Question: Compute the shortest $s \rightarrow t$ path in G .

Assumption: No negative cycles.

- Edges from a node.

Question: What is the max length of the ^{Shortest} path from $s \rightarrow t$

↳ Ans: at most $n-1$.

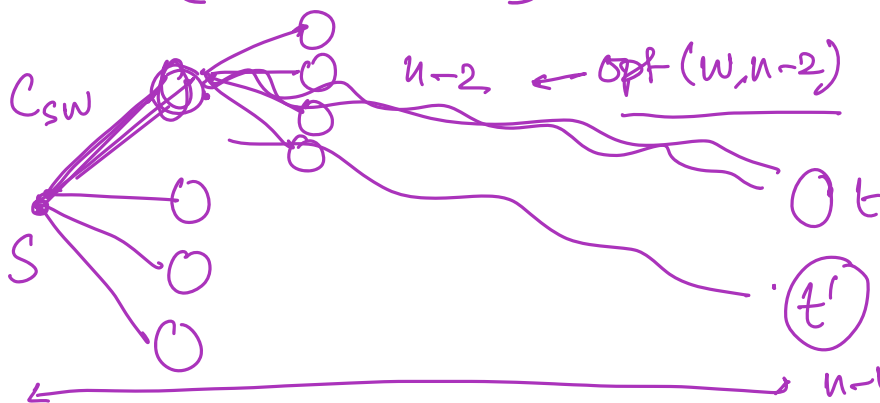
$Opt(v, l)$: min wt of all $v \rightarrow t$ path of length at most l .

$Opt(s, n-1)$: ← This is what we want.

↓
 ↘ Could happen
 paths of length at most $n-2$ ✓

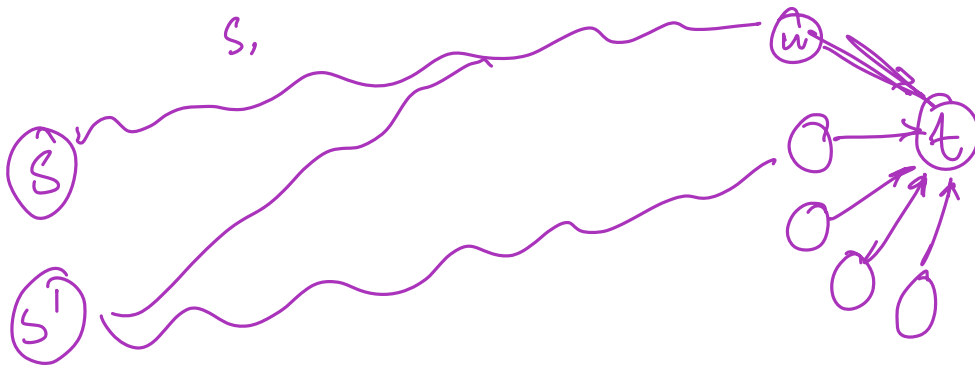
$$\min \{ Opt(w, n-2) + C_{sw} : w \text{ s.t. } (s, w) \in E \}$$

$$Opt(s, n-1) = \min \left\{ \{ Opt(s, n-2) \} \cup \{ Opt(w, n-2) + C_{sw} : w \text{ s.t. } (s, w) \in E \} \right\}$$



$\hat{\text{Opt}}(l, v)$: min wtf over all $s \rightsquigarrow v$ paths of length at most l .

$$\hat{\text{Opt}}(n-1, t) = \min \left\{ \begin{array}{l} \text{Opt}(n-2, t) \\ \{ \text{Opt}(n-2, w) + \underline{C_{wt}} : (w, t) \in E \} \end{array} \right.$$



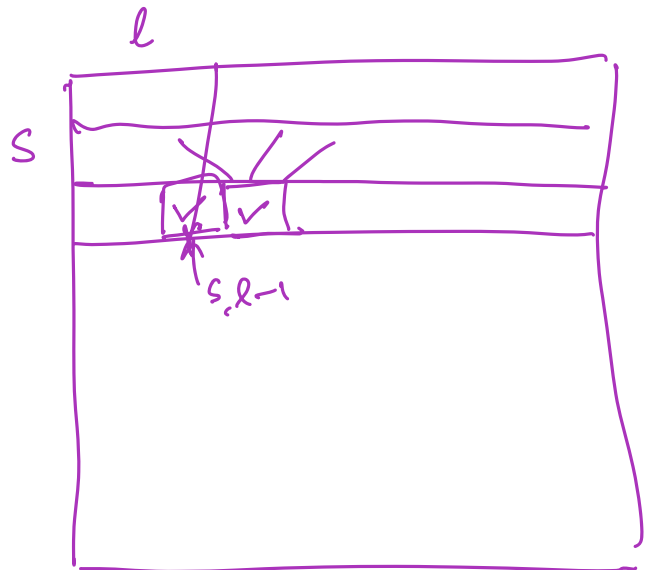
For each entry of the memoization, we consider all the neighbours of that node in consideration.

$$\forall l \in [0, n-1]$$

$$s, l$$

$$w, l-1$$

$$w, l-2$$



$$\underline{O(n^2 + nm)}$$

⑤